Economic Aspects of Information Systems

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The old information systems

I must provide accurate data, expressive interfaces and fast query evaluation to the company's employees

EMP SAL DEPT MNGR

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The new information systems top ten in Google



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Logic databases

Canfly(x,y) :- Canfly(x,z), Flight(z,y). Canfly(x,y) :- Flight(x,y).

- Some of the most elegant work in Database Theory during the 1980s
- Never picked up by practitioners
- Main criticism: "Good for poorly designed databases"

The return of logic

- The Internet is a deliberately badly designed database.
- e.g.:

What is privacy?

- one of society's most vital concerns
 central for e-commerce
 arguably the most crucial and far-reaching current challenge and mission of CS
 least understood scientifically (e.g., is it rational?)
- see, e.g., <u>www.sims.berkeley.edu/~hal</u>, ~/pam,
- [Stanford Law Review, June 2000]

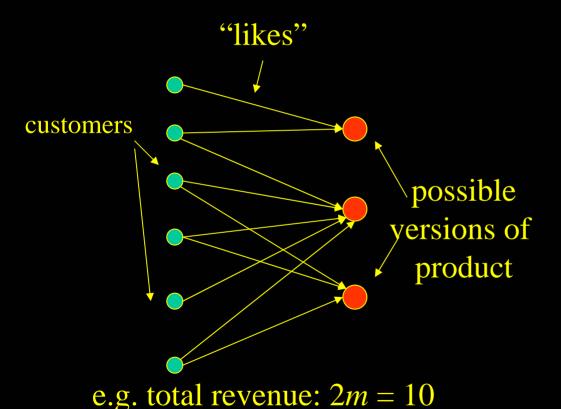
some thoughts on privacy

- also an economic problem
- surrendering private information is either good or bad for you
- example: privacy vs. search costs in computer purchasing

thoughts on privacy (cont.)

- personal information is intellectual property controlled by others, often bearing negative royalty
- selling mailing lists vs. selling aggregate information: false dilemma
- Proposal: Take into account the individual's utility when using personal data for decision-making

e.g., marketing survey



company's utility is proportional to the majority
customer's utility is 1 if in the majority

• how should all participants be compensated?

Collaborative Game Theory

- How should A, B, C split the loot (=20)?
- We are given what each subset can achieve by itself as a function v from the powerset of {A,B,C} to the reals
- $v({}) = 0$

- Values of v
- A: 10
- B: 0
- C: 6
- AB: 14
- BC: 9
- AC: 16
- ABC: 20

first idea (notion of "fairness"): the core

A vector $(x_1, x_2, ..., x_n)$ with $\Sigma_i x_i = v([n])$ (= 20) is in the core if for all *S* we have $x[S] \ge v(S)$

In our example: A gets 11, B gets 3, C gets 6 Problem: Core is often empty (e.g., AB = 15)

second idea: the Shapley value

 $x_i = \mathcal{E}_{\pi}(v[\{j: \pi(j) \le \pi(i)\}] - v[\{j: \pi(j) < \pi(i)\}])$

(Meaning: Assume that the agents arrive at random. Pay each one his/her contribution. Average over all possible orders of arrival.)

Theorem [Shapley]: The Shapley value is the only allocation that satisfies Shapley's axioms.

In our example...

- A gets: 10/3 + 14/6 + 10/6 + 11/3 = 11
- B gets:

0/3 + 4/6 + 3/6 + 4/3 = 2.5

- C gets the rest = 6.5
- NB: Split the cost of a trip among hosts...

- Values of v
- A: 10
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e.g., the UN security council

- 5 permanent, 10 non-permanent
- A resolution passes if voted by a majority of the 15, including all 5 P
- v[S] = 1 if |S| > 7 and S contains 1,2,3,4,5;
 otherwise 0
- What is the Shapley value (~power) of each P member? Of each NP member?

e.g., the UN security council

- What is the probability, when you are the 8th arrival, that all of 1,...,5 have arrived?
- Ans: Choose(10,2)/Choose(15,7) ~ .7%
 Permanent members: ~ 18%

Therefore, $P \neq NP$

third idea: bargaining set fourth idea: nucleolus

seventeenth idea: the von Neumann-Morgenstern solution

0

[Deng and P. 1990] complexity-theoretic critique of solution concepts

Applying to the market survey problem

- Suppose largest *minority* is r
- An allocation is in the core as long as losers get 0, vendor gets > 2r, winners split an amount up to twice their victory margin
- (plus another technical condition saying that split must not be too skewed)

market survey problem: Shapley value

- Suppose margin of victory is at least $\varepsilon > 0\%$
- (realistic, close elections never happen in real life...)
- Vendor gets $m(1+\varepsilon)$
- Winners get $1 + \varepsilon$
- Losers get ε
- (and so, no compensation is necessary)

e.g., recommendation system

- Each participant *i* knows a set of items B_i
- Each benefits 1 from every new item
- Core: empty, unless the sets are disjoint!
- Shapley value: For each item you know, you are owed an amount equal to 1 / (#people who know about it)

--i.e., novelty pays

e.g., collaborative filtering

- Each participant likes/dislikes a set of items (participant is a vector of 0, ±1)
- The "similarity" of two agents is the inner product of their vectors
- There are k "well separated types" (vectors of ± 1), and each agent is a *random perturbation* and *random masking* of a type

collaborative filtering (cont.)

- An agent gets advice on a 0 by asking the most similar other agent who has a ± 1 in that position
- Value of this advice is the product of the agent's true value and the advice.
- How should agents be compensated (or charged) for their participation?

collaborative filtering (result)

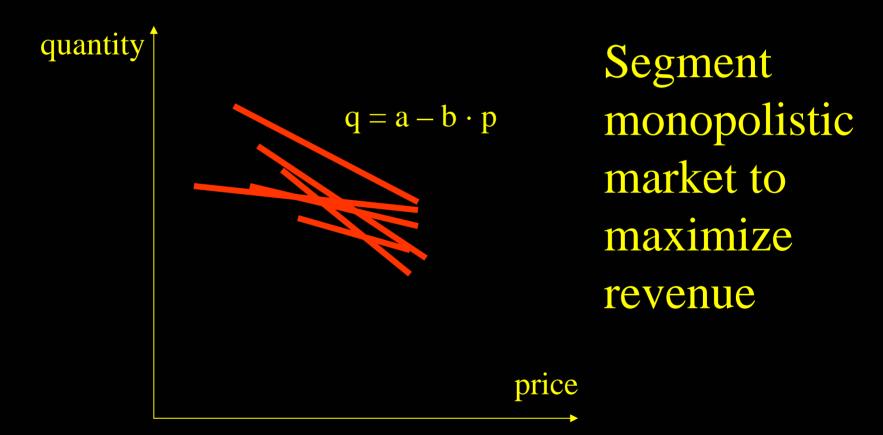
Theorem: An agent's compensation (= value to the community) *is an increasing function of how typical* (close to his/her type) *the agent is.*

The economics of clustering

• The practice of clustering: Confusion, too many criteria and heuristics, no guidelines

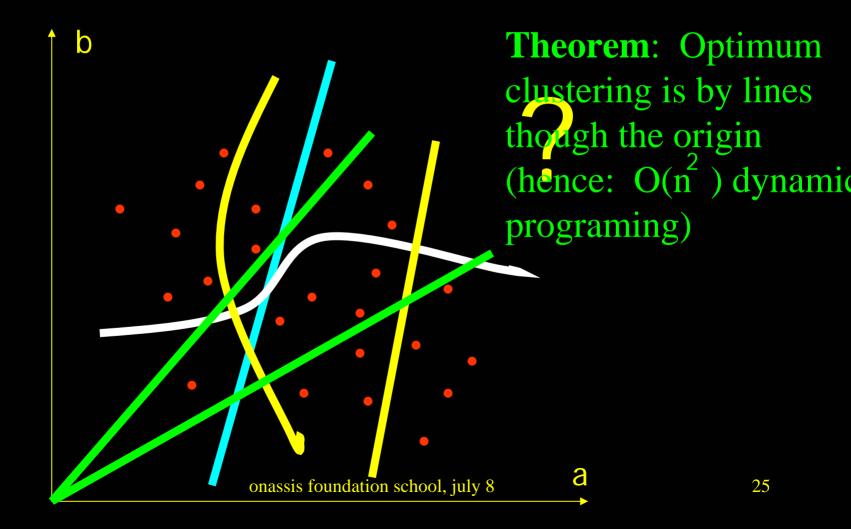
- The theory of clustering: ditto!
- "It's the economy, stupid!" [Kleinberg, P., Raghavan STOC 98, JDKD 99]

Example: market segmentation



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or, in the a – b plane:



So...

- The new databases are selfish agents, and you understand them through economics
- Privacy has an interesting and central economic aspect
- Which gives rise to neat math/algorithmic problems
- And clustering is a meaningful problem only in a well-defined economic context